Towards a Lightweight RDMA Para-Virtualization for HPC

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Background

- Virtualized HPC and HPC Cloud become more popular and requires lightweight solutions (unikernel-like OS)
 - O Cloud: Flexibility, Scalability, Reliability, ...
 - O HPC: Power and extreme speed of computation and data analysis
- Challenging workloads regarding latency throughput, may benefit from RDMA
 - $_{\mbox{O}}$ $\,$ RDMA virtualization is required



Why not SR-IOV ?



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RDMA virtualization: State of the art

Control Path: RDMA verbs translated to host

CENTER

Data Path: RDMA memory mapped to backend driver and RDMA HCA



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virtio-rdma: Design Overview

- High bandwidth, low latency network I/O
- Lightweight, flexible, and portable
- Targeting at unikernel-like OS, **OSv**:
 - O Fast, lightweight, less overhead
 - O Bridge Cloud and HPC
- Support socket and RDMA verbs API Host
- Shared memory for the intra-host communication
- Support RoCE* and InfiniBand

*RoCE: RDMA over Converged Ethernet



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APP

Socket / RDMA verbs

Guest



VisorHPC @ HiPEAC 2017



Guest

APP

Socket / RDMA verbs

virtio-rdma: Control Path

- Re-use the HyV backend driver
 - O Ported to a newer Linux kernel

Guest side implementation on Osv

Re-implemented frontend driver,
 hypercall

O Streamlined/minimized OFED support

- Data path: RDMA memory is mapped directly from guest to host and HCA
 - Completion Queue, Queue Pairs, and
 Work Requests



Lightweight, flexible and portable



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virtio-rdma: Context Switch

Context switches are reduced

- O OSv runs application in the same privilege level, single address space
- O No boundary of user/kernel space
- O Less calls to external libraries
- O minimum number of memory copy
 - O only a few memory copies are needed for OFED core data structure



virtio-rdma: Memory Mapping

- Guest allocate contiguous memory in virtual address space
- frontend driver remap the virtual address to guest physical address with chunks
 - HyV needs to take care of the noncontiguous physical address and the offset in the page
 - OSv always tries to allocate contiguous mem.
 - virtio-rdma only cares for the offset
- Backend driver remaps the chunks into

host physical address Efficiency on address translation:







virtio-rdma: Future Work

Shared memory

- O Based on IVSHMEM
- O Protocol switch in RDMA verbs API
- O Communication buffers, e.g. user memory, shared to the VMs of same host
- O Update RDMA memory regions directly, e.g. Completion Queue

Socket API support

- $_{\mbox{O}}$ forward to use RDMA and shared memory communications
- O re-implement rsocket or libvma onto OSv
- O no modification to the user application is needed

HPC Integration

O Setup the necessary environment via Torque when VM starts



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Demo

root@host1:/home/demo



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Conclusion

Redesigned and Implemented the new virtio-rdma driver on OSv

- O Compatible with HyV backend driver, i.e. same hypercall
- O Fewer OFED dependencies and fewer context switches
- O Simpler memory translation
- O Flexible and portable for other unikernel-like OS
- Streamlined OFED user libraries and core structures for OSV
- Reused the HyV backend driver and ported it to newer Linux kernel
- HPC Cloud enabled





Acknowledgement

- This work is supported by the MIKELANGELO project co-founded by the European Commission (H2O2O framework programme)
- Many thanks to the co-authors
- Special thanks to Jonas Pfefferle at IBM Zurich (author of HyV)





Thank you!